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# ISyE8813

## Alternating direction method of multipliers (ADMM)

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Consider the problem

$$\begin{aligned} \min & f(x) + g(x) \\ \text{s.t.} & Ax + By = b \in \mathbb{R}^r \end{aligned}$$

where  $f \in \overline{\text{Conv}}(\mathbb{R}^n)$  and  $g \in \overline{\text{Conv}}(\mathbb{R}^m)$ . For a given  $\rho > 0$ , define

$$L_\rho(x, y; p) = f(x) + g(x) + p^T(Ax + By - b) + \frac{\rho}{2}\|Ax + By - b\|^2$$

## 1 Augmented Lagrangian Method

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**Algorithm 1** (ALM Method)

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0. Let  $p_0 \in \mathbb{R}^r$  be given.

1. Compute

$$(x, y) \in \text{Argmin}_{(x', y')} L_\rho(x', y'; p_0). \quad (1)$$

2. Set

$$p = p_0 + \rho(Ax + By - b)$$

3. Set  $p_0 \leftarrow p$  and go to step 1.

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**Optimality Conditions for (1)**

$$0 \in \partial f(x) + A^*[p_0 + \rho(Ax + By - b)]$$

$$0 \in \partial g(y) + B^*[p_0 + \rho(Ax + By - b)]$$

$$\frac{p_0 - p}{\rho} = b - Ax - By$$

or equivalently,

$$0 \in \partial f(x) + A^*p$$

$$0 \in \partial g(y) + B^*p$$

$$\frac{p_0 - p}{\rho} = b - Ax - By \quad (2)$$

We will now see that ALM is an exact implementation of PPM for a certain Bregman distance. Indeed, for any  $z = (x, y; p) \in \mathbb{R}^n \times \mathbb{R}^m \times \mathbb{R}^r$ , define

$$T(z) = T(x, y; p) := \begin{bmatrix} 0 & 0 & A^* \\ 0 & 0 & B^* \\ -A & -B & 0 \end{bmatrix} \begin{bmatrix} x \\ y \\ p \end{bmatrix} + \begin{bmatrix} \partial f(x) \\ \partial g(y) \\ b \end{bmatrix}.$$

and the distance generating function

$$w(z) = w(x, y; p) := \frac{1}{2} \|p\|^2$$

Then, it can be shown that  $T$  is maximal monotone and that

$$(dw)_{z_0}(z) = \frac{1}{2} \|p - p_0\|^2 \tag{3}$$

and

$$\nabla w(z) = \begin{bmatrix} 0 \\ 0 \\ p \end{bmatrix}$$

Also, (18) becomes

$$\frac{\nabla w(z_0) - \nabla w(z)}{\rho} \in T(z)$$

This is the equation of an exact PPM with respect to the Bregman distance (3) and prox stepsize  $\lambda$  equal to  $\rho$

## 2 Generalized HPE framework

Consider the MIP

$$0 \in T(z)$$

where  $T : \mathbb{R}^n \rightrightarrows \mathbb{R}^n$  is maximal monotone.

Assume that for some semi-norm in  $\mathbb{R}^n$ , the following conditions hold:

- 1)  $T^{-1}(0) \neq \emptyset$
- 2) there exists  $m, M > 0$  such that for every  $z, z' \in \mathbb{R}^n$ , have

$$\begin{aligned} (dw)_z(z') &\geq \frac{m}{2} \|z' - z\|^2 \\ \|\nabla w(z') - \nabla w(z)\|^* &\leq M \|z' - z\| \end{aligned}$$

where

$$\|\cdot\|^* = \sup\{\langle \cdot, v \rangle : \|v\| \leq 1\}$$

Condition 2) implies that

$$\frac{m}{2} \|z - z'\|^2 \leq (dw)_z(z') \leq \frac{M}{2} \|z - z'\|^2 \quad \forall z, z' \in Z. \quad (4)$$

**Example:** If

$$w(\cdot) = (1/2) \|\cdot\|_Q^2$$

where  $Q$  is a self-adjoint positive semidefinite linear operator, then  $w$  satisfies condition 2 with  $(m, M) = (1, 1)$ .

**Proposition 2.1.** *If  $Q : \mathbb{R}^n \rightarrow \mathbb{R}^n$  is a self-adjoint positive semidefinite linear operator, then the semi-norm*

$$\|\cdot\| := \langle \cdot, Q(\cdot) \rangle^{1/2}$$

*satisfies the following statements:*

- (a)  $\text{dom } \|\cdot\|^* = \text{Im}(Q)$  and

$$\|Qz\|^* = \|z\| \quad \forall z \in \mathbb{R}^n$$

- (b) if  $Q$  is invertible, then

$$\|z\|^* = \langle Q^{-1}z, z \rangle^{1/2} \quad \forall z \in \mathbb{R}^n$$

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**Algorithm 2** (Inexact PPM framework with Bregman distance)

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0. Let  $z_0 \in \mathbb{R}^n$  and  $\sigma \in [0, 1]$  be given.

1. Find  $\lambda > 0$  and  $(z, \tilde{z}, \varepsilon) \in \mathbb{R}^n \times \mathbb{R}^n \times \mathbb{R}_+$  such that

$$\frac{\nabla w(z_0) - \nabla w(z)}{\lambda} \in T^\varepsilon(\tilde{z}) \quad (5)$$

$$(dw)_z(\tilde{z}) + \lambda\varepsilon \leq \sigma(dw)_{z_0}(\tilde{z}) \quad (6)$$

2. Set  $z_0 \leftarrow z$  and go to step 1.

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Hence, sequence-wise, we have for every  $k \geq 1$  that

$$r_k := \frac{\nabla w(z_{k-1}) - \nabla w(z_k)}{\lambda_k} \in T^{\varepsilon_k}(\tilde{z}_k) \quad (7)$$

$$(dw)_z(\tilde{z}_k) + \lambda_k \varepsilon_k \leq \sigma(dw)_{z_{k-1}}(\tilde{z}_k) \quad (8)$$

**Proposition 2.2. (Pointwise)** *Assume that  $\sigma < 1$  and  $\lambda_k \geq \underline{\lambda}$  for every  $k \geq 1$ . Then, for every  $k \geq 1$ , there exists  $i \leq k$  such that*

$$\|r_i\|^* \leq \frac{2M}{\underline{\lambda}\sqrt{mk}} \sqrt{\frac{(1+\sigma)(dw)_0}{1-\sigma}} = \mathcal{O}\left(\frac{1}{\underline{\lambda}\sqrt{k}}\right)$$

and

$$\varepsilon_i \leq \frac{(1+\sigma)(dw)_0}{(1-\sigma)\underline{\lambda}k} = \mathcal{O}\left(\frac{1}{\underline{\lambda}k}\right)$$

where  $r_i$  is as in (7) and

$$(dw)_0 := \inf \{(dw)_{z_0}(z^*) : z^* \in T^{-1}(0)\}$$

For  $k \geq 1$ , define  $\Lambda_k := \sum_{i=1}^k \lambda_i$  and the ergodic iterate  $(\tilde{z}_k^a, r_k^a, \varepsilon_k^a)$  as

$$\tilde{z}_k^a = \frac{1}{\Lambda_k} \sum_{i=1}^k \lambda_i \tilde{z}_i, \quad r_k^a := \frac{1}{\Lambda_k} \sum_{i=1}^k \lambda_i r_i, \quad \varepsilon_k^a := \frac{1}{\Lambda_k} \sum_{i=1}^k \lambda_i (\varepsilon_i + \langle r_i, \tilde{z}_i - \tilde{z}_k^a \rangle). \quad (9)$$

**Theorem 2.3. (Ergodic convergence of the NE-HPE)** *For every  $k \geq 1$ , we have*

$$r_k^a \in T^{\varepsilon_k^a}(\tilde{z}_k^a)$$

and

$$\|r_k^a\|^* \leq \frac{2\sqrt{2}M(dw_0)^{1/2}}{\sqrt{m}\Lambda_k}, \quad \varepsilon_k^a \leq \left(\frac{3M}{m}\right) \left[\frac{3(dw)_0 + \sigma\theta_k}{\Lambda_k}\right].$$

where

$$\theta_k := \max_{i=1, \dots, k} (dw)_{z_{i-1}}(\tilde{z}_i). \quad (10)$$

Moreover, the sequence  $\{\theta_k\}$  is bounded under either one of the following situations:

(a)  $\sigma < 1$ , in which case

$$\theta_k \leq \frac{(dw)_0}{1 - \sigma}; \quad (11)$$

(b)  $\text{Dom } T$  is bounded, in which case

$$\theta_k \leq \frac{2M}{m} [(dw)_0 + D]$$

where

$$D := \sup \{ \min \{ (dw)_y(y'), (dw)_{y'}(y) \} : y, y' \in \text{Dom } T \}$$

### 3 ADDM

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**Algorithm 3** (ADMM)

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0. Let  $(y_0, p_0) \in \mathbb{R}^m \times \mathbb{R}^r$  be given.

1. Compute

$$x \in \text{Argmin}_{x'} L_\rho(x', y_0; p_0). \quad (12)$$

2. Compute

$$y \in \text{Argmin}_{y'} L_\rho(x, y'; p_0). \quad (13)$$

3. Set

$$p = p_0 + \rho(Ax + By - b)$$

4. Set  $(y_0, p_0) \leftarrow (y, p)$  and go to step 1.

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**Optimality Conditions for (17) and (18)**

$$0 \in \partial f(x) + A^*[p_0 + \rho(Ax + By_0 - b)]$$

$$0 \in \partial g(y) + B^*[p_0 + \rho(Ax + By - b)]$$

$$\frac{p_0 - p}{\rho} = b - Ax - By$$

or equivalently,

$$0 \in \partial f(x) + A^*\tilde{p}$$

$$0 \in \partial g(y) + B^*p \quad (14)$$

$$\frac{p_0 - p}{\rho} = b - Ax - By$$

where

$$\tilde{p} = p_0 + \rho(Ax + By_0 - b)$$

Hence,

$$p - \tilde{p} = \rho B(y - y_0) \quad (15)$$

and so the second inclusion in (14) implies that

$$\begin{aligned} \rho B^* B(y_0 - y) &\in \partial g(y) + B^*p + \rho B^* B(y_0 - y) \\ &= \partial g(y) + B^*[p + B(y_0 - y)] = \partial g(y) + B^*\tilde{p} \end{aligned}$$

Thus, (14) is equivalent to

$$0 \in \partial f(x) + A^*\tilde{p}$$

$$\rho B^* B(y_0 - y) \in \partial g(y) + B^*\tilde{p} \quad (16)$$

$$\frac{p_0 - p}{\rho} = b - Ax - By$$

Defining

$$\tilde{z} := (x, y; \tilde{p}), \quad w(z) := \frac{1}{2\rho} \|p\|^2 + \frac{\rho}{2} \|By\|^2$$

and noting that

$$\nabla w(z) = \begin{bmatrix} 0 \\ \rho B^* B y \\ \rho^{-1} p \end{bmatrix}$$

it follows that (16) is equivalent to

$$\nabla w(z_0) - \nabla w(z) \in T(\tilde{z})$$

which shows that (5) is satisfied with  $\varepsilon = 0$  and  $\lambda = 1$ .

We will now show that (6) is also satisfied with  $\varepsilon = 0$ ,  $\lambda = 1$  and  $\sigma = 1$ . Indeed, since  $\varepsilon = 0$ , the left hand side of (6) becomes

$$(dw)_z(\tilde{z}) + \lambda\varepsilon = (dw)_z(\tilde{z}) = \frac{1}{2\rho} \|\tilde{p} - p\|^2 = \frac{\rho}{2} \|B(y - y_0)\|^2$$

where the last equality is due to (15). On the other hand, since  $\sigma = 1$ , the right hand side of (6) becomes

$$\sigma(dw)_{z_0}(\tilde{z}) = (dw)_{z_0}(\tilde{z}) = \frac{1}{2\rho} \|\tilde{p} - p_0\|^2 + \frac{\rho}{2} \|B(y - y_0)\|^2$$

Thus, (6) holds with  $\varepsilon = 0$ ,  $\lambda = 1$  and  $\sigma = 1$ .

Is the quantity  $\theta_k$  in (10) bounded? Since  $\sigma = 1$ , this fact is not clear. Note that

$$\begin{aligned} (dw)_{z_0}(\tilde{z}) &= \frac{1}{2\rho} \|\tilde{p} - p_0\|^2 + \frac{\rho}{2} \|B(y - y_0)\|^2 \\ &\leq \frac{1}{\rho} [\|\tilde{p} - p\|^2 + \|p - p_0\|^2] + \frac{\rho}{2} \|B(y - y_0)\|^2 \\ &= \frac{1}{\rho} [\rho^2 \|B(y - y_0)\|^2 + \|p - p_0\|^2] + \frac{\rho}{2} \|B(y - y_0)\|^2 \\ &= \frac{1}{\rho} \|p - p_0\|^2 + \frac{3\rho}{2} \|B(y - y_0)\|^2 \leq 3(dw)_{z_0}(z) \\ &\leq 6[(dw)_{z_0}(z^*) + (dw)_z(z^*)] \end{aligned}$$

Since the quantity  $(dw)_z(z^*)$  remains bounded (why?), it follows that  $(dw)_{z_0}(\tilde{z})$  also remains bounded, and so does the quantity  $\theta_k$  in (10)

## 4 Proximal ADMM

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### Algorithm 4 (Proximal ADMM)

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0. Let  $\theta > 0$ ,  $(y_0, p_0) \in \mathbb{R}^m \times \mathbb{R}^r$ , and positive-semidefinite self-adjoint linear operators  $\mathcal{G}$  and  $\mathcal{H}$  be given.

1. Compute

$$x \in \operatorname{Argmin}_{x'} L_\rho(x', y_0; p_0) + \frac{1}{2} \|x - x_{k-1}\|_{\mathcal{G}}^2. \quad (17)$$

2. Compute

$$y \in \operatorname{Argmin}_{y'} L_\rho(x, y'; p_0) + \frac{1}{2} \|y - y_{k-1}\|_{\mathcal{H}}^2 \quad (18)$$

3. Set

$$p = p_0 + \theta\rho(Ax + By - b)$$

4. Set  $(y_0, p_0) \leftarrow (y, p)$  and go to step 1.

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**Theorem 4.1. (Pointwise convergence of the proximal ADMM)** *Consider the sequence  $\{(x_k, y_k, p_k)\}$  generated by the proximal ADMM with  $\theta \in (0, (\sqrt{5} + 1)/2)$ , and let  $\{\tilde{x}_k\}$  be defined as*

$$\tilde{p}_k = p_{k-1} + \rho(Ax_k + By_{k-1} - b). \quad (19)$$

*Then, for every  $k \in \mathbb{N}$ ,*

$$\begin{pmatrix} G(x_{k-1} - x_k) \\ (H + \rho B^* B)(y_{k-1} - y_k) \\ (p_{k-1} - p_k)/(\rho\theta) \end{pmatrix} \in \begin{bmatrix} \partial f(x_k) + A^* \tilde{p}_k \\ \partial g(y_k) + B^* \tilde{p}_k \\ Ax_k + By_k - b \end{bmatrix} \quad (20)$$

*and there exists  $i \leq k$  such that*

$$\left( \|x_{i-1} - x_i\|_{\mathcal{G}}^2 + \|y_{i-1} - y_i\|_{(\mathcal{H} + \rho B^* B)}^2 + \frac{1}{\beta\theta} \|p_{i-1} - p_i\|_{\mathcal{X}}^2 \right) = \mathcal{O}\left(\frac{1}{k}\right)$$

## 5 Chambolle and Poch algorithm

Consider the problem

$$(P) \quad \min_x \max_y \langle Kx, y \rangle + G(x) - F^*(y)$$

where  $G \in \overline{\text{Conv}}(\mathbb{R}^n)$ ,  $F \in \overline{\text{Conv}}(\mathbb{R}^m)$ ,  $K : \mathbb{R}^n \mapsto \mathbb{R}^m$  is linear. The problem (P) is equivalent to

$$\min_x F(Kx) + G(x)$$

and has the dual formulation

$$\max_x \min_y \langle Kx, y \rangle + G(x) - F^*(y) =: \psi(x, y)$$

or equivalently

$$\max_y -G^*(-K^*y) - F^*(y).$$

Furthermore, assume that there exists  $(x^*, y^*) \in \mathbb{R}^n \times \mathbb{R}^m$  such that

$$-Kx^* + \partial F^*(y^*) \ni 0, \quad K^*y^* + \partial G(x^*) \ni 0$$

or equivalently

$$(0, 0) \in \partial [\psi(\cdot, y^*) - \psi(x^*, \cdot)](x^*, y^*)$$

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### Algorithm 5 (Chambolle-Poch (CP) Algorithm)

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0. Choose  $\tau_1, \tau_2 > 0$ ,  $\theta = 1$ ,  $(p_0, q_0) \in \mathbb{R}^n \times \mathbb{R}^m$  and set  $\bar{p}_0 = p_0$  and  $k = 1$
1. compute

$$\begin{aligned} q_k &= (I + \tau_2 \partial F^*)^{-1}(q_{k-1} + \tau_2 K \bar{p}_{k-1}) \\ p_k &= (I + \tau_1 \partial G)^{-1}(p_{k-1} - \tau_1 K^* q_k) \\ \bar{p}_k &= p_k + \theta(p_k - p_{k-1}) \end{aligned}$$

2. set  $k \leftarrow k + 1$  and go to step 1.
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We have

$$\frac{q_k - q_{k-1}}{\tau_2} - K \bar{p}_{k-1} + \partial F^*(q_k) \ni 0 \tag{21}$$

$$\frac{p_k - p_{k-1}}{\tau_1} + K^* q_k + \partial G(p_k) \ni 0 \tag{22}$$

**Proposition 5.1.** *The CP algorithm is an instance of the inexact IPP framework with  $\sigma^2 := \|K\|^2 \tau_1 \tau_2$ .*

Fix  $k \geq 1$  and let

$$\tilde{x}_k = p_k, \quad x_k = \theta p_k, \quad \tilde{y}_k = q_k$$

Then, (22) implies that

$$\frac{x_{k-1} - x_k}{\theta \tau_1} \in K^* \tilde{y}_k + \partial G(\tilde{x}_k)$$

and (21) implies that

$$\begin{aligned}
0 &\in \frac{q_k - q_{k-1}}{\tau_2} - K\bar{p}_{k-1} + \partial F^*(q_k) \\
&= \frac{q_k - q_{k-1}}{\tau_2} + K(\bar{p}_k - \bar{p}_{k-1}) + K(p_k - \bar{p}_k) - Kp_k + \partial F^*(q_k) \\
&= \frac{q_k - q_{k-1}}{\tau_2} + K(\bar{p}_k - \bar{p}_{k-1}) - \theta K(p_k - p_{k-1}) - Kp_k + \partial F^*(q_k) \\
&= \frac{y_k - y_{k-1}}{\tau_2} - K\tilde{x}_k + \partial F^*(\tilde{y}_k)
\end{aligned}$$

where

$$y_k := q_k + \tau_2 K[\bar{p}_k - \theta p_k]$$

Now, consider the distance generating function

$$w(x, y) = \frac{1}{\tau_1} \|x\|^2 + \frac{1}{\tau_2} \|y\|^2$$

Then, if  $z = (x_k, y_k)$ ,  $\tilde{z} = (\tilde{x}_k, \tilde{y}_k)$  and  $z_0 = (x_{k-1}, y_{k-1})$ , then

$$\begin{aligned}
(dw)_z(\tilde{z}) &= \frac{1}{\tau_1} \|x_k - \tilde{x}_k\|^2 + \frac{1}{\tau_2} \|y_k - \tilde{y}_k\|^2 = \frac{1}{\tau_1} \|\theta p_k - p_k\|^2 + \frac{1}{\tau_2} \|y_k - q_k\|^2 \\
&= \frac{(1-\theta)^2}{\tau_1} \|p_k\|^2 + \frac{1}{\tau_2} \|\tau_2 K(\bar{p}_k - \theta p_k)\|^2 = \frac{(1-\theta)^2}{\tau_1} \|p_k\|^2 + \tau_2 \|K\|^2 \|\bar{p}_k - \theta p_k\|^2 \\
&= \frac{(1-\theta)^2}{\tau_1} \|p_k\|^2 + \tau_2 \|K\|^2 \|p_k - \theta p_{k-1}\|^2
\end{aligned}$$

Also,

$$\begin{aligned}
(dw)_{z_0}(\tilde{z}) &= \frac{1}{\tau_1} \|x_{k-1} - \tilde{x}_k\|^2 + \frac{1}{\tau_2} \|y_{k-1} - \tilde{y}_k\|^2 \\
&\geq \frac{1}{\tau_1} \|x_{k-1} - \tilde{x}_k\|^2 = \frac{1}{\tau_1} \|\theta p_{k-1} - p_k\|^2
\end{aligned}$$

## 6 A block-decomposition algorithm

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### Algorithm 6 (Block-Decomposition Algorithm)

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0. Choose  $\tau_1, \tau_2 > 0$ ,  $(p_0, q_0) \in \mathbb{R}^n \times \mathbb{R}^m$  and set  $k = 1$
1. Compute

$$\begin{aligned}\tilde{q}_k &= (I + \tau_2 \partial F^*)^{-1}(q_{k-1} + \tau_2 K p_{k-1}) \\ p_k &= (I + \tau_1 \partial G)^{-1}(p_{k-1} - \tau_1 K^* \tilde{q}_k) \\ q_k &= \tilde{q}_k + \tau_2 K [p_k - p_{k-1}]\end{aligned}$$

2. Set  $k \leftarrow k + 1$  and go to step 1
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Have

$$\frac{\tilde{q}_k - q_{k-1}}{\tau_2} - K p_{k-1} + \partial F^*(\tilde{q}_k) \ni 0 \quad (23)$$

$$\frac{p_k - p_{k-1}}{\tau_1} + K^* \tilde{q}_k + \partial G(p_k) \ni 0 \quad (24)$$

Fix  $k \geq 1$  and let

$$\tilde{x}_k = x_k = p_k, \quad \tilde{y}_k = \tilde{q}_k, \quad y_k = q_k$$

Then, (24) implies that

$$\frac{x_{k-1} - x_k}{\tau_1} \in K^* \tilde{y}_k + \partial G(\tilde{x}_k)$$

and (23) implies that

$$\begin{aligned}0 &\in \frac{\tilde{q}_k - q_{k-1}}{\tau_2} - K p_{k-1} + \partial F^*(\tilde{q}_k) \\ &= \frac{\tilde{q}_k - q_{k-1}}{\tau_2} + K(p_k - p_{k-1}) - K p_k + \partial F^*(\tilde{y}_k) \\ &= \frac{\tilde{q}_k - q_{k-1}}{\tau_2} + \frac{q_k - \tilde{q}_k}{\tau_2} - K p_k + \partial F^*(\tilde{y}_k) \\ &= \frac{q_k - q_{k-1}}{\tau_2} - K p_k + \partial F^*(\tilde{y}_k) \\ &= \frac{y_k - y_{k-1}}{\tau_2} - K \tilde{x}_k + \partial F^*(\tilde{y}_k)\end{aligned}$$

and hence that

$$\frac{y_{k-1} - y_k}{\tau_2} \in -K \tilde{x}_k + \partial F^*(\tilde{y}_k)$$

Now, consider the distance generating function

$$w(x, y) = \frac{1}{\tau_1} \|x\|^2 + \frac{1}{\tau_2} \|y\|^2$$

Then, if  $z = (x_k, y_k)$ ,  $\tilde{z} = (\tilde{x}_k, \tilde{y}_k)$  and  $z_0 = (x_{k-1}, y_{k-1})$ , then

$$\begin{aligned} (dw)_z(\tilde{z}) &= \frac{1}{\tau_1} \|x_k - \tilde{x}_k\|^2 + \frac{1}{\tau_2} \|y_k - \tilde{y}_k\|^2 = \frac{1}{\tau_2} \|q_k - \tilde{q}_k\|^2 \\ &= \frac{1}{\tau_2} \|\tau_2 K(p_k - p_{k-1})\|^2 \leq \tau_2 \|K\|^2 \|p_k - p_{k-1}\|^2 \end{aligned}$$

Also,

$$\begin{aligned} (dw)_{z_0}(\tilde{z}) &= \frac{1}{\tau_1} \|x_{k-1} - \tilde{x}_k\|^2 + \frac{1}{\tau_2} \|y_{k-1} - \tilde{y}_k\|^2 \\ &\geq \frac{1}{\tau_1} \|x_{k-1} - \tilde{x}_k\|^2 = \frac{1}{\tau_1} \|p_{k-1} - p_k\|^2 \end{aligned}$$

Hence, if

$$\sigma := \tau_1 \tau_2 \|K\|^2 \leq 1$$

then

$$(dw)_z(\tilde{z}) \leq \tau_2 \|K\|^2 \|p_k - p_{k-1}\|^2 = \frac{\sigma}{\tau_1} \|p_k - p_{k-1}\|^2 \leq \sigma (dw)_{z_0}(\tilde{z})$$

Hence, the two conditions for the IPP framework holds with  $\varepsilon = 0$  and  $\lambda = 1$